Mitosis - Transparently Self-Replicating Page-Tables for Large-Memory Machines



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Mitosis: Mitigate NUMA Effects on Page Table Walks



Up to 3.2x speed up for single threaded applications.



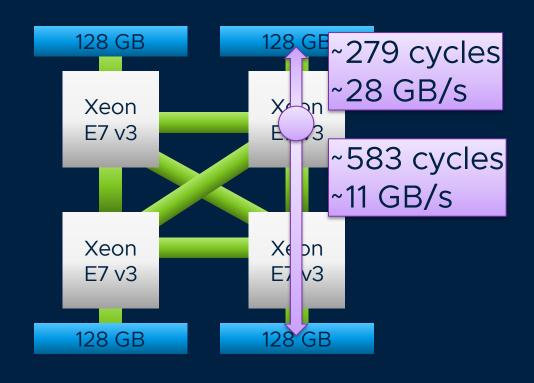
Up to 1.3x speed up on multithreaded applications.



No modifications of the application workloads.



NUMA Effects on Large Multi-Socket Machines



Measured on 4x12 Intel Xeon E7-4850 v3, with 512GB RAM (4x128GB)

Bandwidth & capacity limited per processor socket

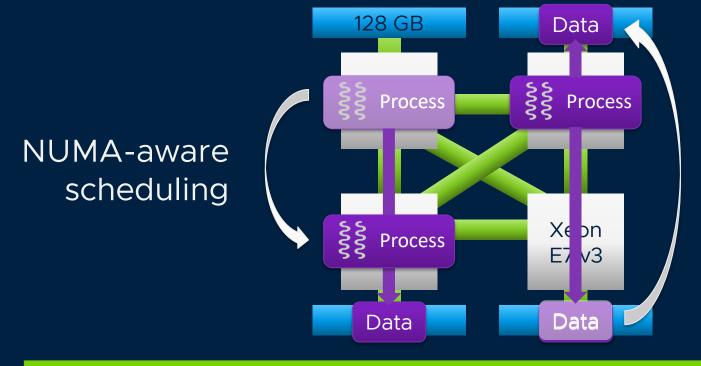
More bandwidth & capacity needed → Multi-socket machine

Non-uniform memory access (NUMA) machine



Handling NUMA Effects on Data Pages

Goal: Keep process threads and data close



NUMA-aware data migration

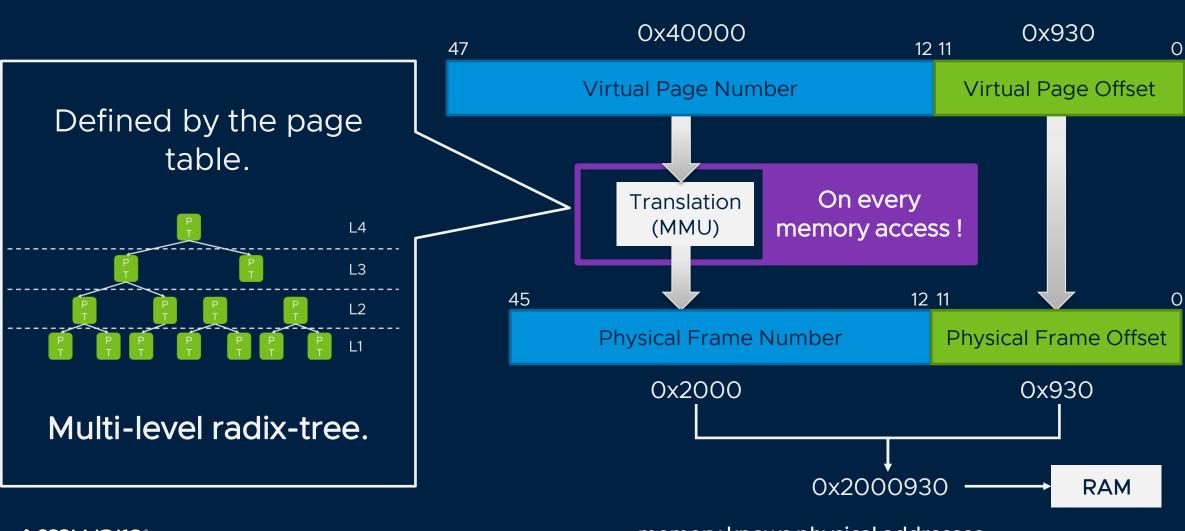
Decades of work has focused on data placement.

What about placement of page tables?



Every Memory Access Requires Virtual-to-Physical Translation

Background (x86 specific)

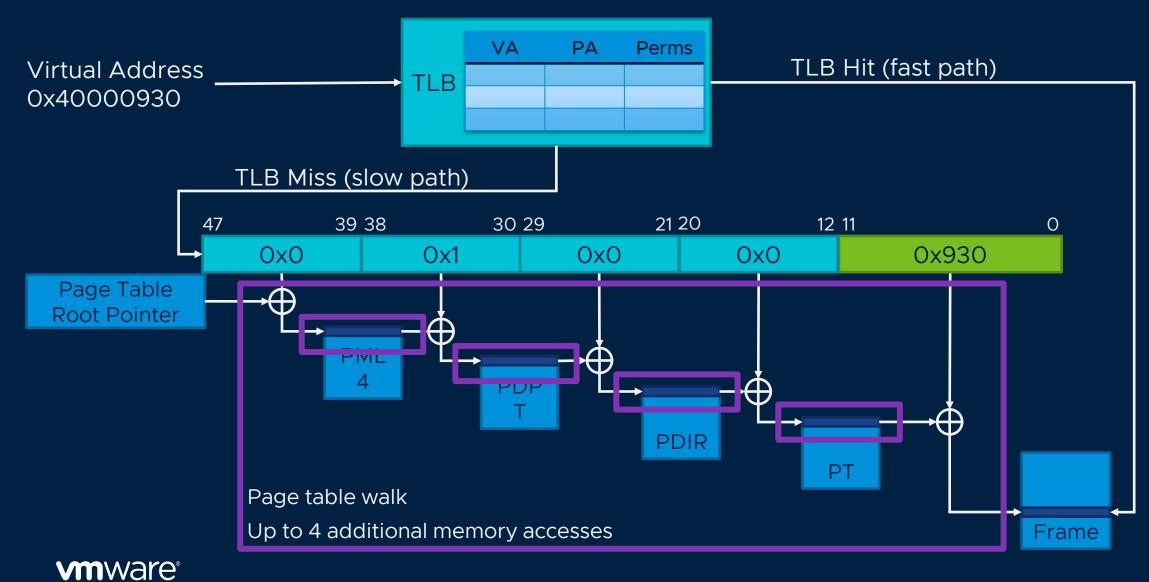




Processes deal with virtual addresses

Translation Lookaside Buffer

TLB misses trigger page table walks

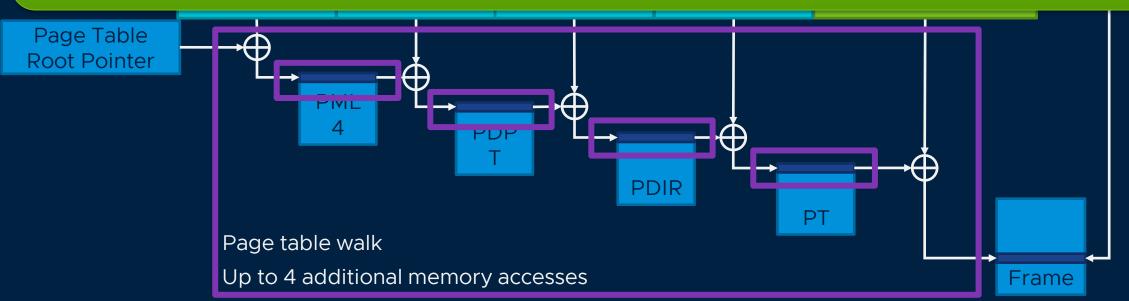


Translation Lookaside Buffer

TLB misses trigger page table walks

Key Problem:

- TLB coverage << available memory → more frequent TLB misses.
 Test machine with 512GB RAM, TLB coverage < 1% of memory
- 2. These TLB misses require multiple memory accesses





NUMA effects on page table walks?





Study: Page Table Walks and NUMA

How often is the page table remote and what's the resulting slowdown?

Multi-Threaded Scenario

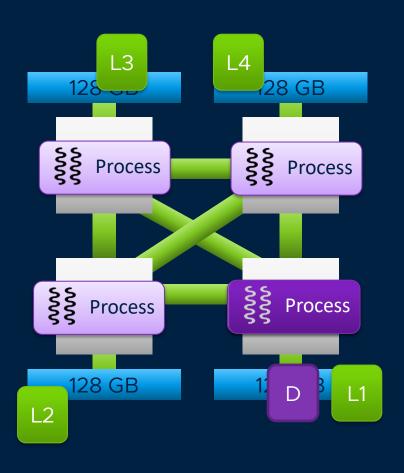
Distribution of page table entries

- Local NUMA Node
- Remote NUMA Node

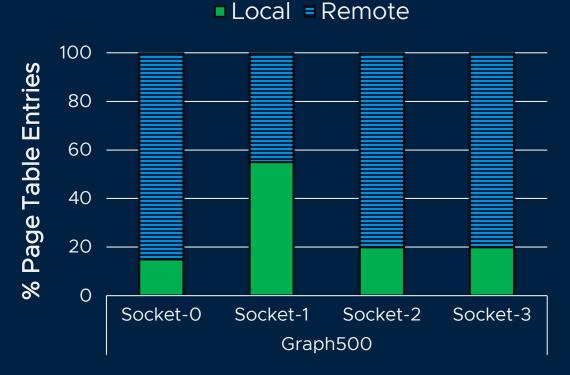


Many Threads – One Page Table

Remote Memory Accesses are Inevitable



Majority of TLB misses require remote memory accesses



Measured on Linux v4.17



Study: Page Table Walks and NUMA

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Multi-Threaded Scenario

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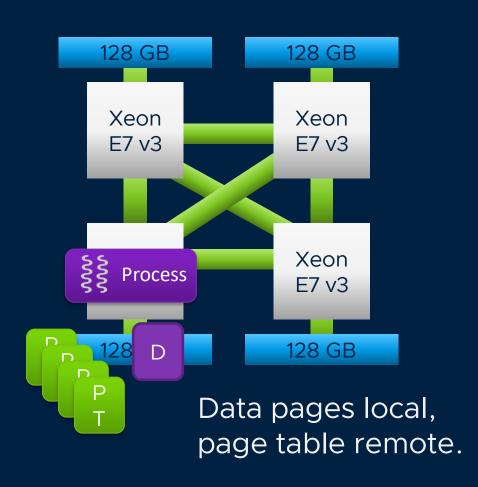
Workload-Migration Scenario

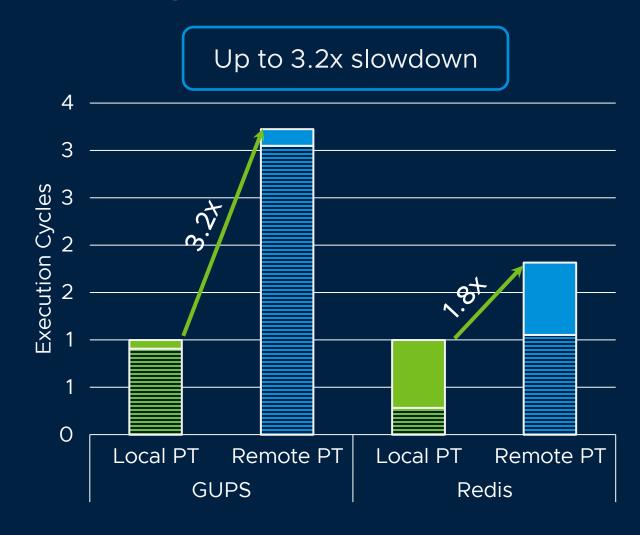
Slow down due to remote memory accesses during page table walks



Remote Page Walk Slow Down

Data Migration Policy moves Data Pages, but not Page Tables









Study: Page Table Walks and NUMA

How often is the page table remote and what's the resulting slowdown?

Multi-Threaded Scenario

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Majority of TLB misses require remote memory accesses

Workload-Migration Scenario

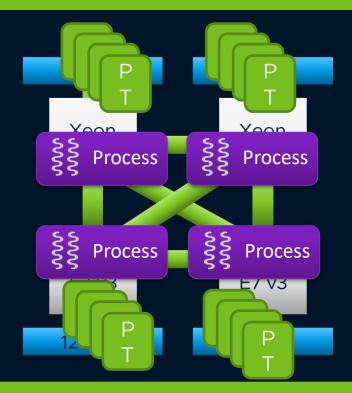
Slow down due to remote memory accesses during page table walks

Up to 3.2x slowdown



Mitosis: Transparent Replication and Migration for Page-Tables Approach

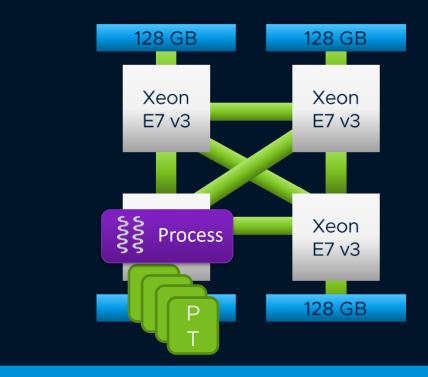
Multi-Threaded Scenario



Mitosis **replicates** the page table

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Workload-Migration Scenario



Mitosis migrates the page table

Mitosis: Keep Page Tables Local using Replication and Migration.

Implementation on Linux / x86_64

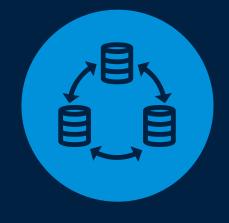




Mechanism: Replication of Page Tables



Replica Allocation



Consistent updates

(Correctness)



Replica Utilization

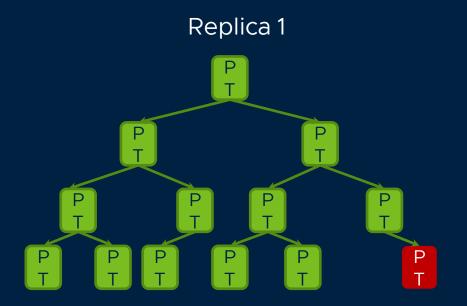
(Performance)

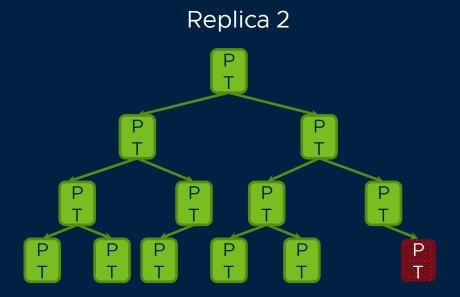


Allocation of Page-Table Replicas



Mitosis eagerly allocates memory for each page table replica.



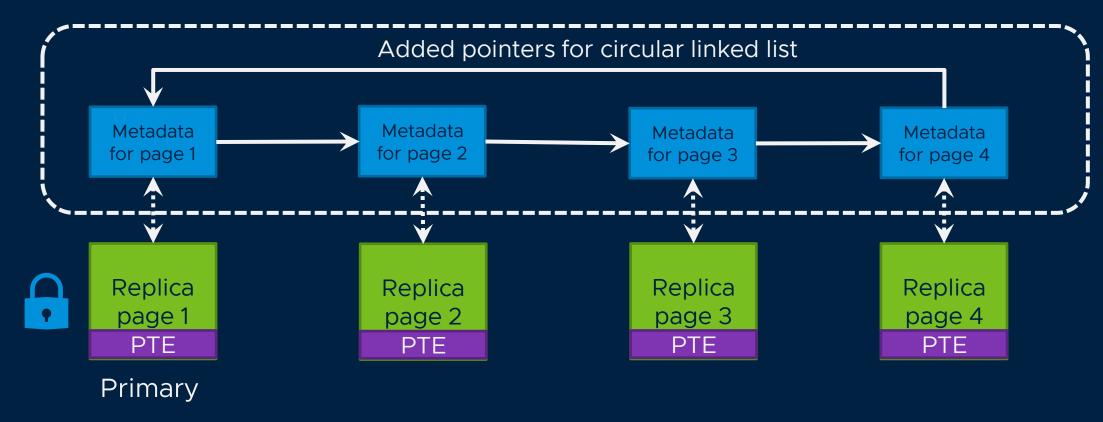




Consistent Updates to Page-Table Replicas Correctness



Mitosis eagerly keeps replicas consistent on update





Utilizing Page Table Replicas

Performance



Scheduler picks thread to run

Context switch

Select local replica

Write translation base register

Extend Linux' mm_struct
with an array of page table roots
Indexed by NUMA node

mm->root[local_node()]



PV-Ops: Paravirtualization Operations in the Linux Kernel

PV-Ops Interface

- Set page table entry
- Allocate and free page table page.
- Reading / writing translation base register

native bare metal environment
 set_pte() → native_set_pte()

paravirtualized environments
 set_pte() → xen_set_pte()

Mitosis PV-Ops Backend

Mitosis implemented as PV-Ops backend.

Extends the native functions to propagate updates to replicas

Adding functions to read entries

-> access & dirty bits



Consistency: Dealing with Access and Dirty Bits

Set by hardware and read by the OS

6 (D) Dirty; indicates whether software has written to the 4-KByte page referenced by this entry (see Section 4.8)	5 (A)	Accessed; indicates whether software has accessed the 4-KByte page referenced by this entry (see Section 4.8)
γ γ γ γ γ γ γ γ γ γ γ γ γ γ γ γ γ γ γ	6 (D)	Dirty; indicates whether software has written to the 4-KByte page referenced by this entry (see Section 4.8)

Bits set by hardware only in local replica > Inconsistency!

Correctness Problem:
OS relies on those bits in paging / caching decisions

```
if (pte_young(*ptep) || pte_dirty(*ptep)) {
  // do something
}
```

Solution: OR-ing entries of all replicas to determine dirty / accessed state!



Mitosis Policy Settings

System Wide Policy

procfs entries to globally enable/disable Mitosis

/proc/sys/kernel/pgtable_replication

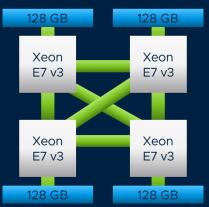
Per Application Policy

Integration in numactl / libnuma

numactl -r 0-3 <workload>



Results

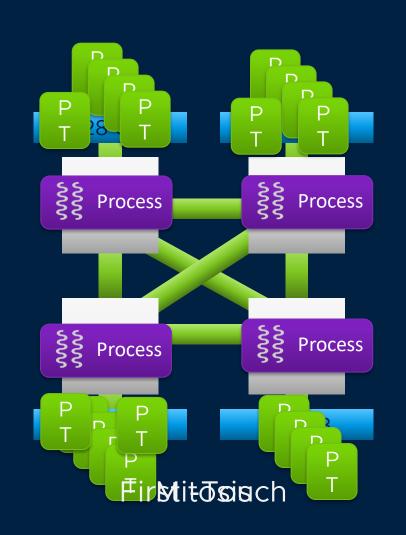


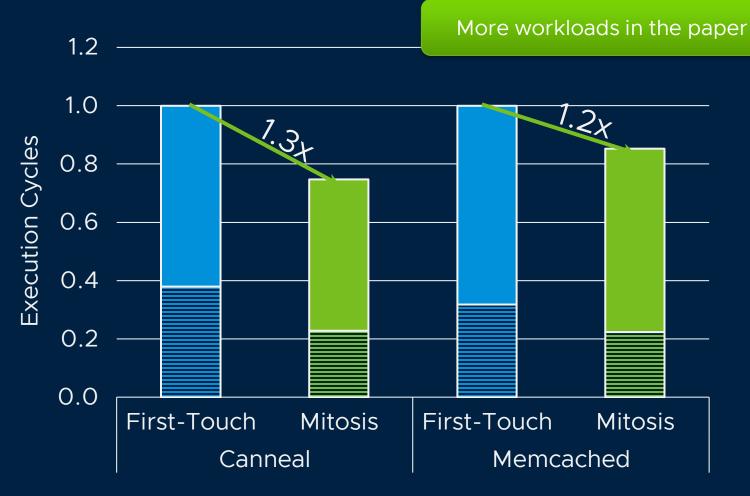
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Evaluation: Multi-Threaded Scenario with Mitosis



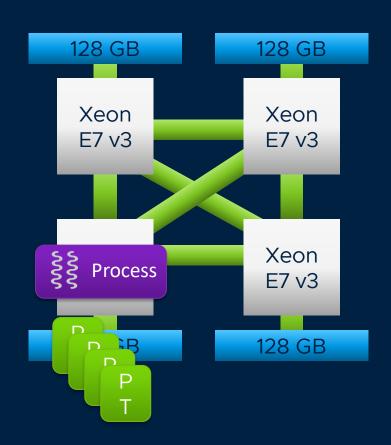


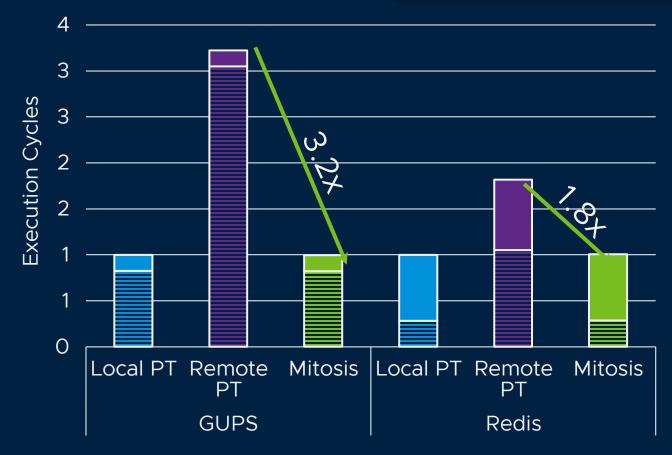
Performance improvement with replication of page tables



Evaluation: Workload-Migration Scenario with Mitosis

More workloads in the paper



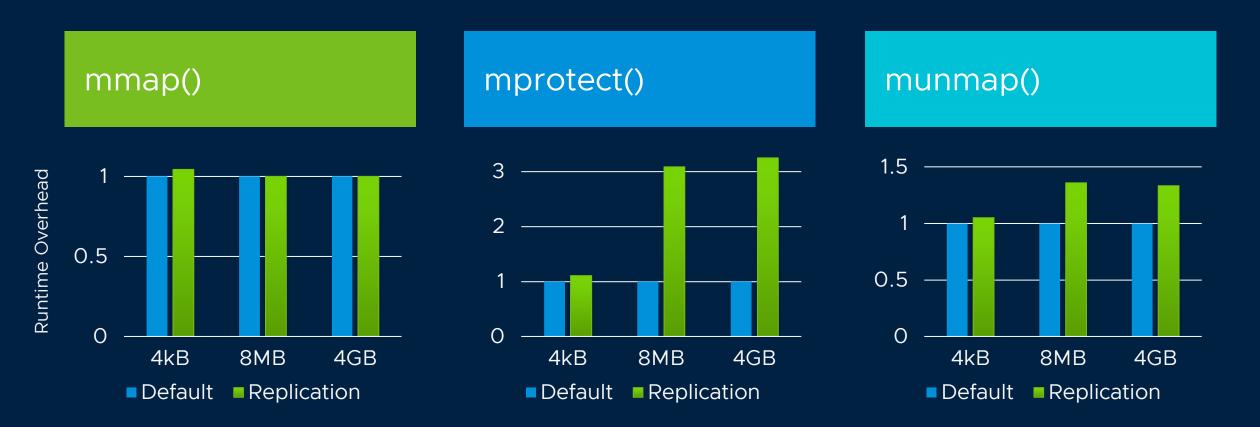


Performance improvement with migration of page tables



Replica Maintenance Overhead

Mitosis with 4-Way Replication



Space Overhead: 0.6% of additional memory for 4-way replication



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Mitosis mitigates NUMA effects on page table walks

NUMA effects

The first study to show NUMA effects on page table walks.

Mitosis Linux **Implementation**

Page table replication as PV-Ops backend.

Integration into libnuma and numactl.

Results

3.2x speed up for workload migration scenario

Up to 1.3x speed up for multi-socket scenario.

