



Smelt: Machine-aware Atomic Broadcast Trees for Multicores

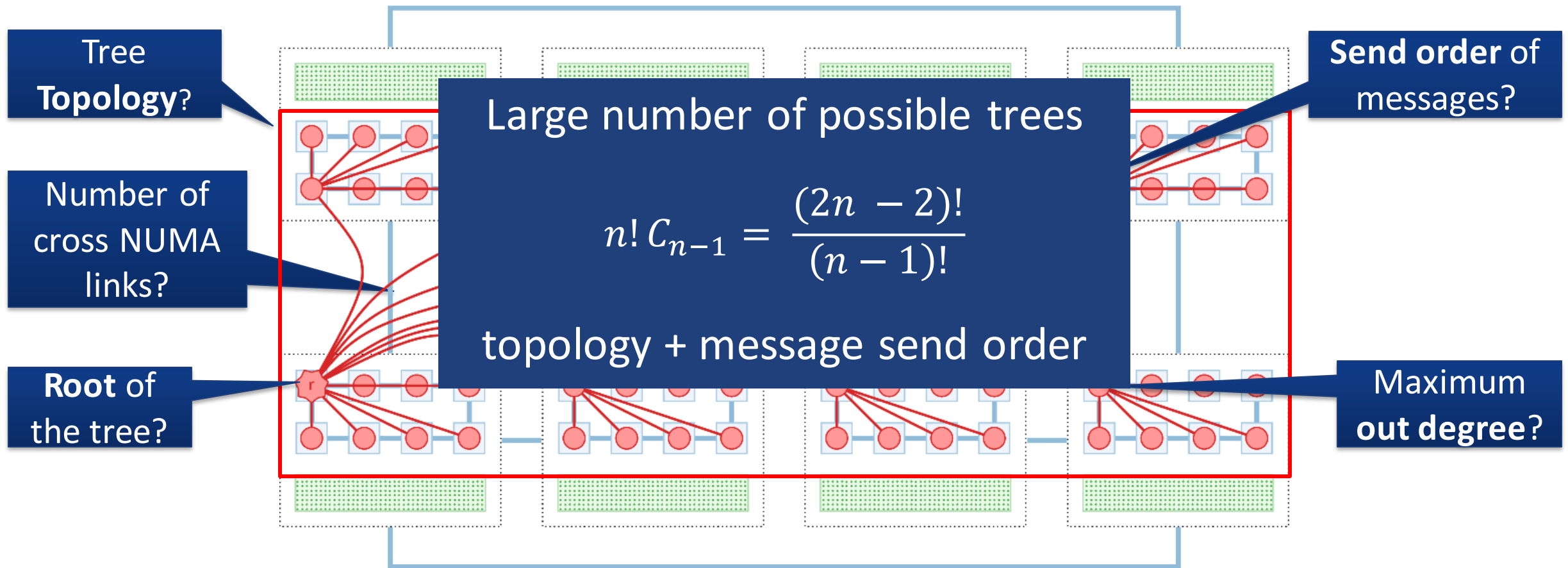
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Large number of trees: topologies and send orders

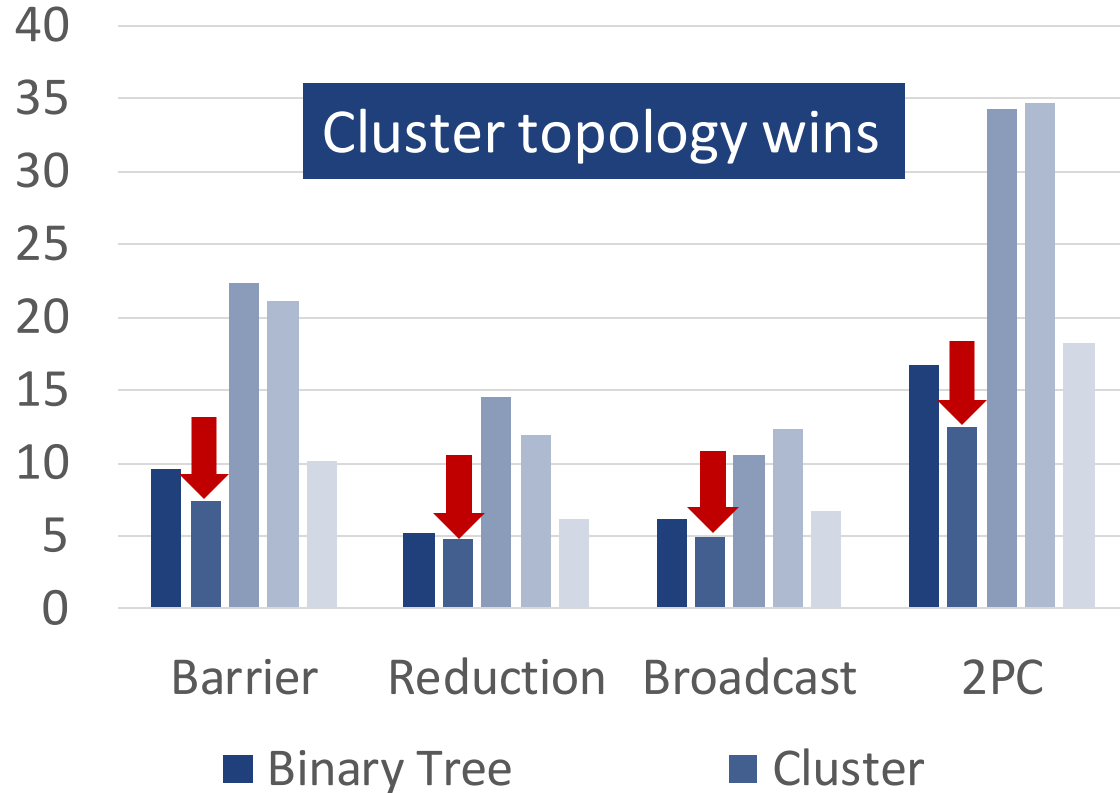
■ CPU
 ■ L3 Cache
 socket
— interconnect
— tree topology



There is no globally optimal tree structure

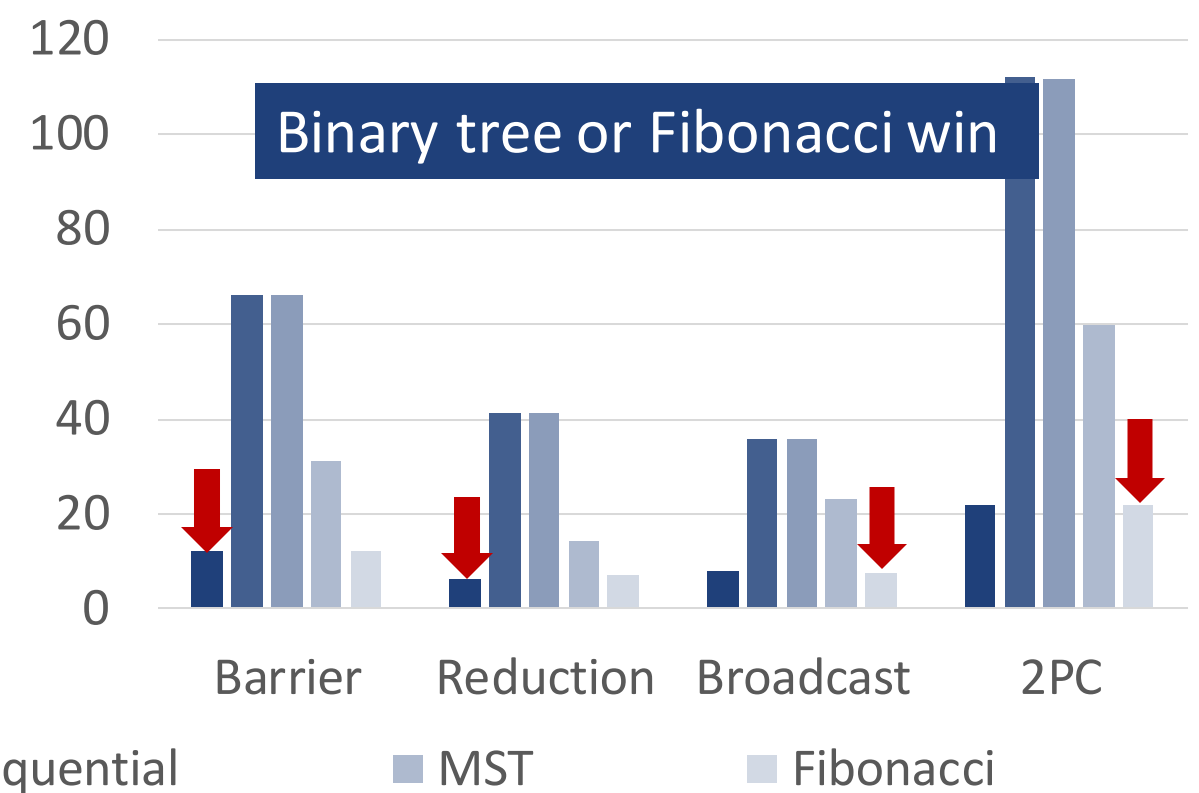
AMD Interlagos (4 Socket x 4 Cores x 2 Threads)

Execution Time [kCycles]

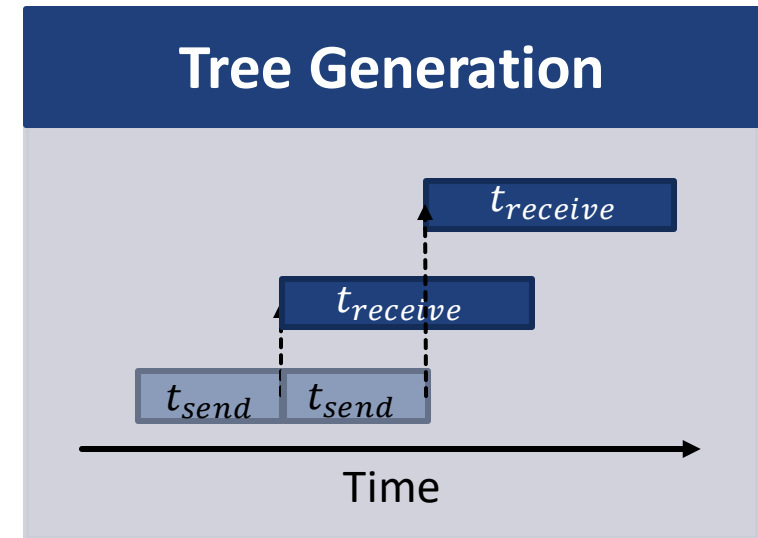
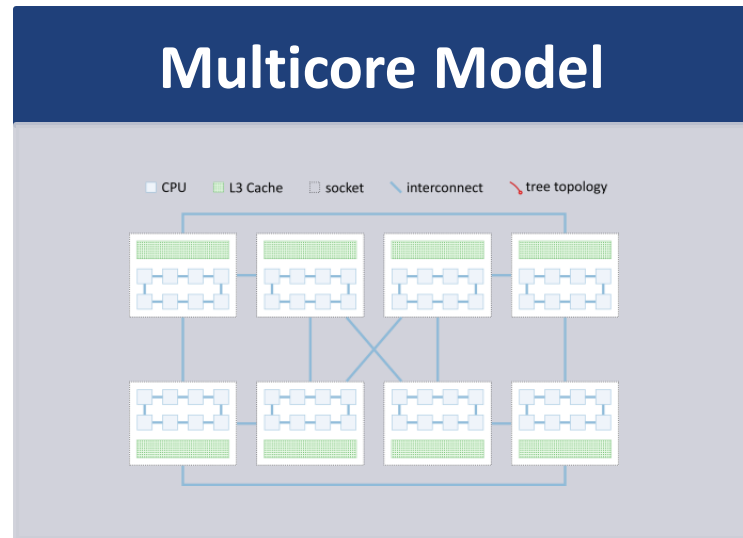


Intel Xeon Phi (61 Cores x 1 Thread)

Execution Time [kCycles]



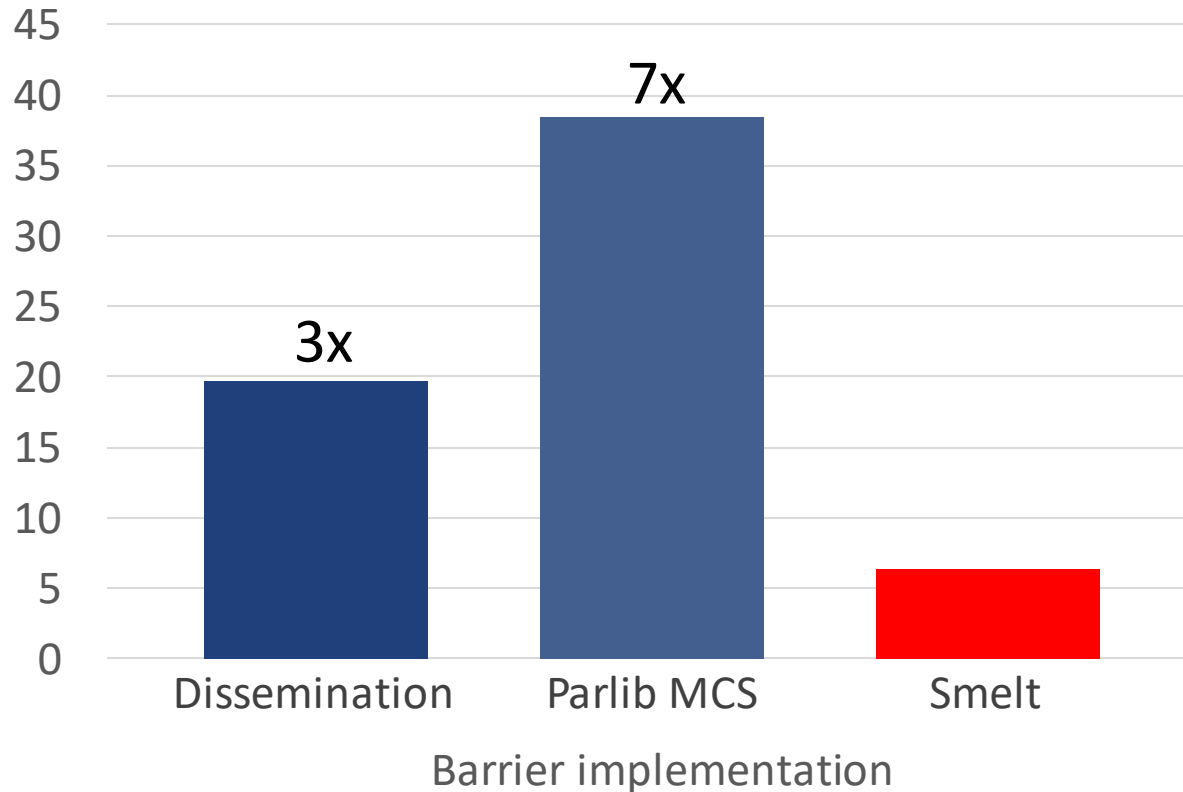
Smelt: Automatic optimization of broadcast and reduction trees



Example: Building fast and simple barriers

Barrier Benchmark on Intel Sandy Bridge 4x8x2

Execution Time [kCycles]



Dramatic improvement through
automatic optimization of
communication patterns

Broadcasts and reductions are central building blocks for parallel programs

Performance

Atomic broadcasts

Replication for **data locality**

e.g. Shoal, Carrefour, SMMP OS, FOS

Fault-Tolerance

Agreement protocols, atomic broadcasts

Replication for **failure resilience**

e.g. 1Paxos

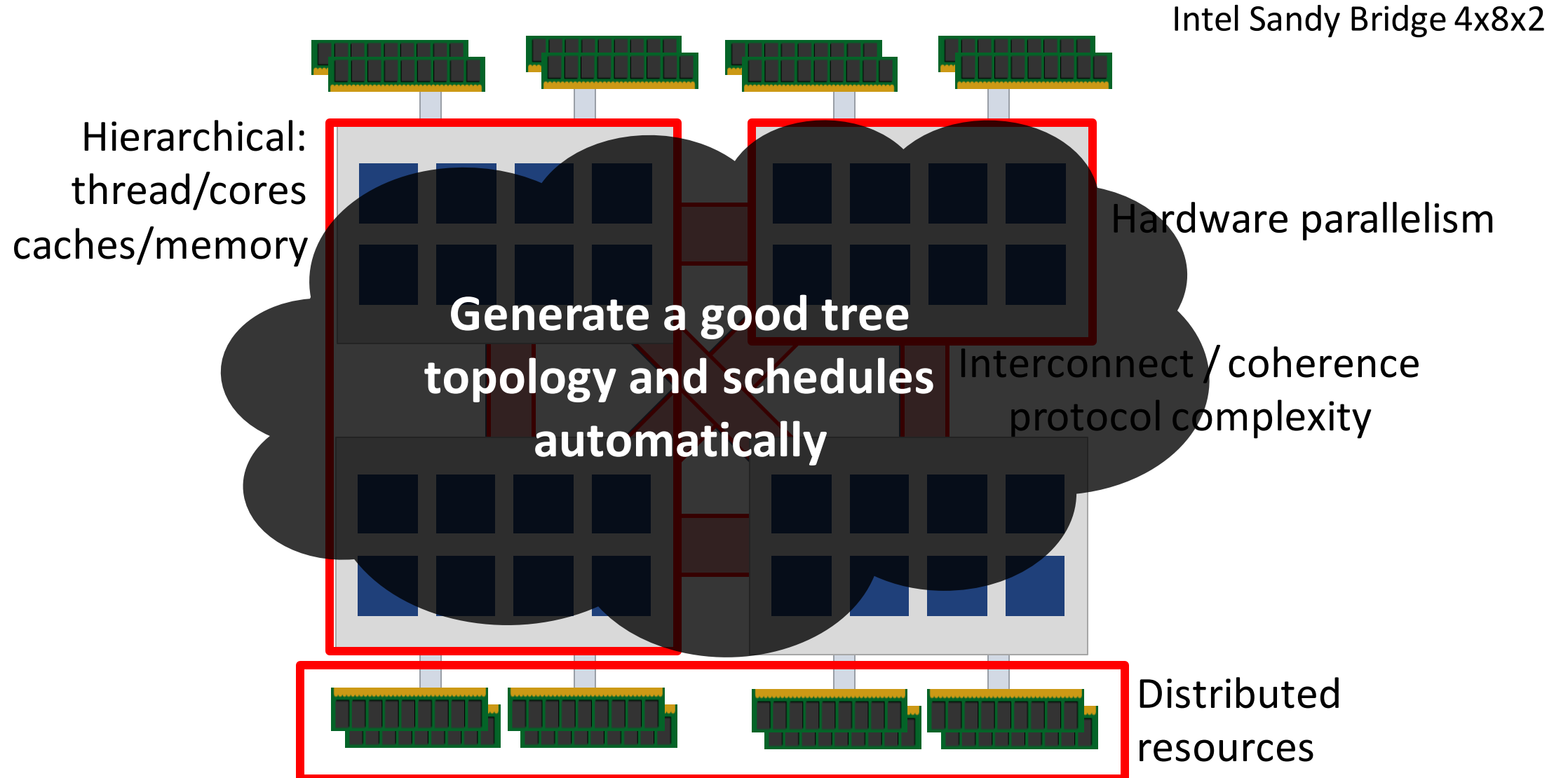
Execution Control

Reductions, broadcast, barriers

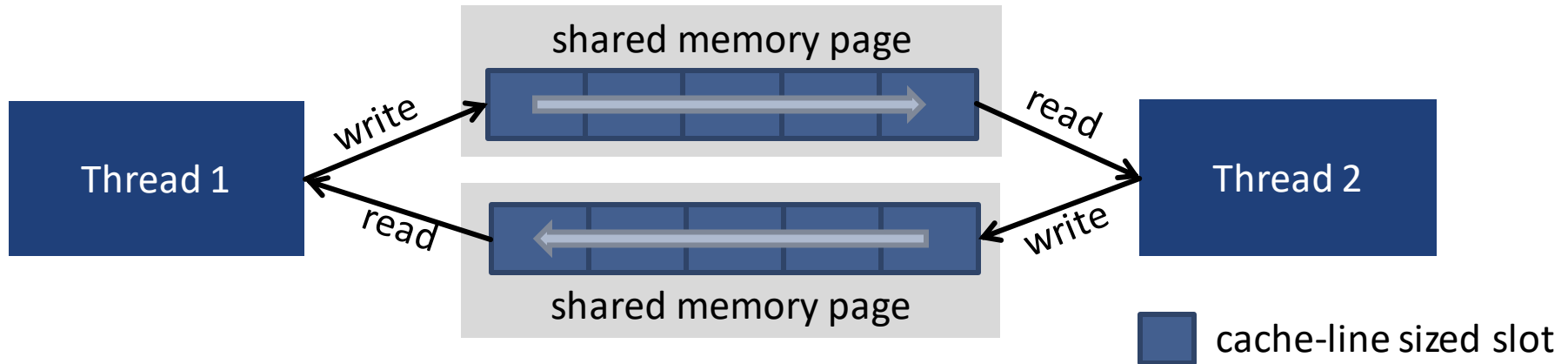
Thread **synchronization**, data gathering

e.g. OpenMP

Multicore hardware is complex



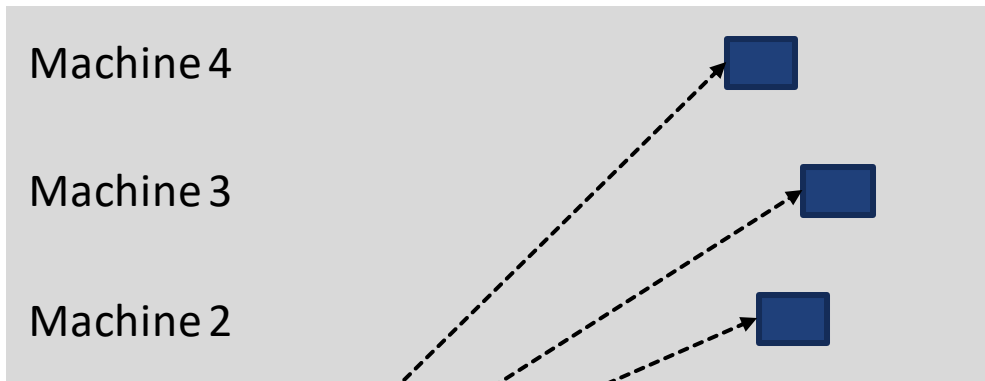
Smelt is based on peer-to-peer message passing



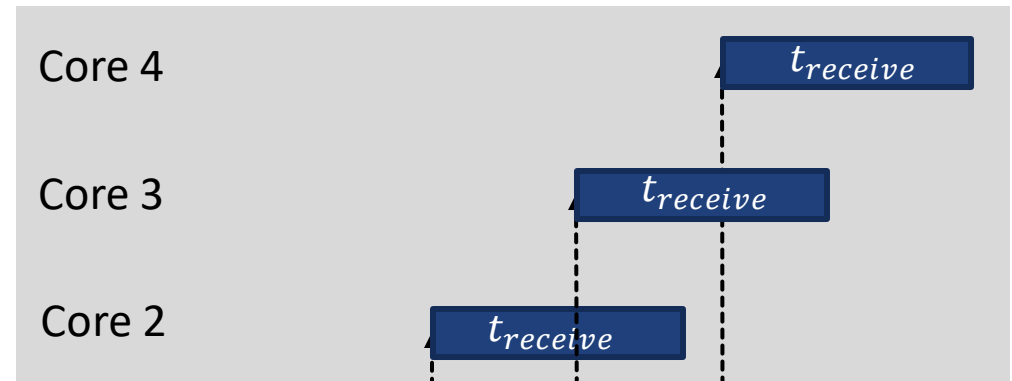
- Works well for our approach.
- Clear concept: Enables **reasoning** about send and receive costs

Message-passing on multicores is different

Classical Network



Multicore interconnect

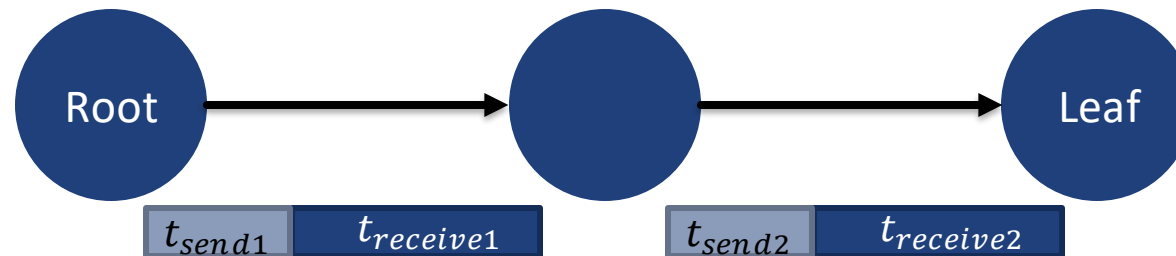


On multicores send and receive times dominate propagation time

Goal: Minimize total time of the broadcast

Minimizing the total time of a broadcast

- $t_{broadcast} = t_{last} - t_{start}$
- Minimize the longest path from the root to the leaves.



$$t_{path} = \sum (t_{send} + t_{receive})$$

We need to know the send and receive cost between any pair of cores

Information obtained from hardware discovery

AMD Interlagos 4x4x2

```
$ lscpu
```

NUMA distance: abstract value

Doesn't distinguish between
send() and recv()

Symmetric: $A \rightarrow B \Rightarrow B \rightarrow A$

```

L1d cache: 16K
L1i cache: 16K
L2 cache: 2048K
L3 cache: 6144K
NUMA node0 CPU(s): 0,4,8,12,16,20,24,28
NUMA node1 CPU(s): 32,36,40,44,48,52,56,60
NUMA node2 CPU(s): 2,6,10,14,18,22,26,30
NUMA node3 CPU(s): 34,38,42,46,50,54,58,62
NUMA node4 CPU(s): 3,7,11,15,19,23,27,31
NUMA node5 CPU(s): 35,39,43,47,51,55,59,63
NUMA node6 CPU(s): 1,5,9,13,17,21,25,29
NUMA node7 CPU(s): 33,37,41,45,49,53,57,61

```

```
$ numactl -hardware
```

node distances:

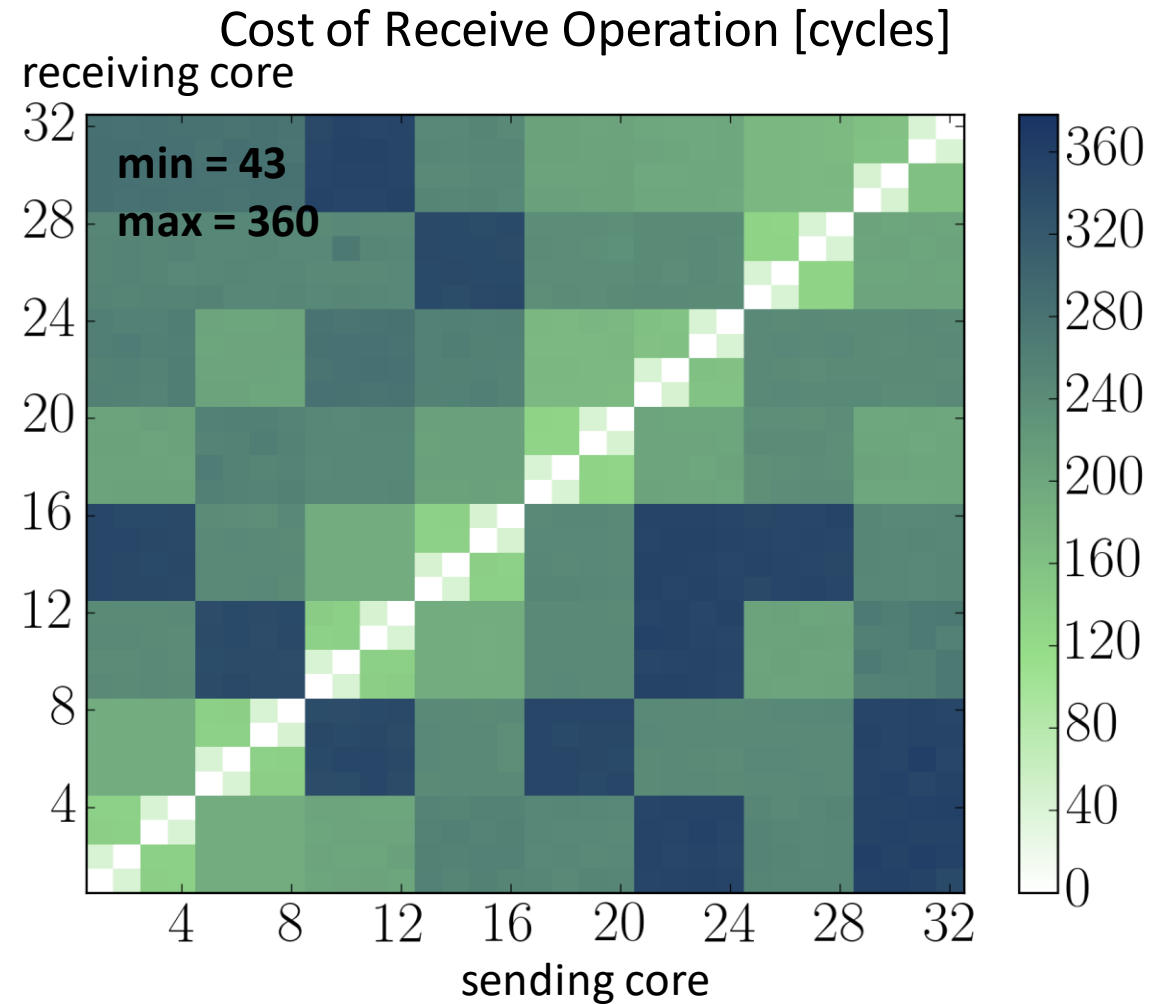
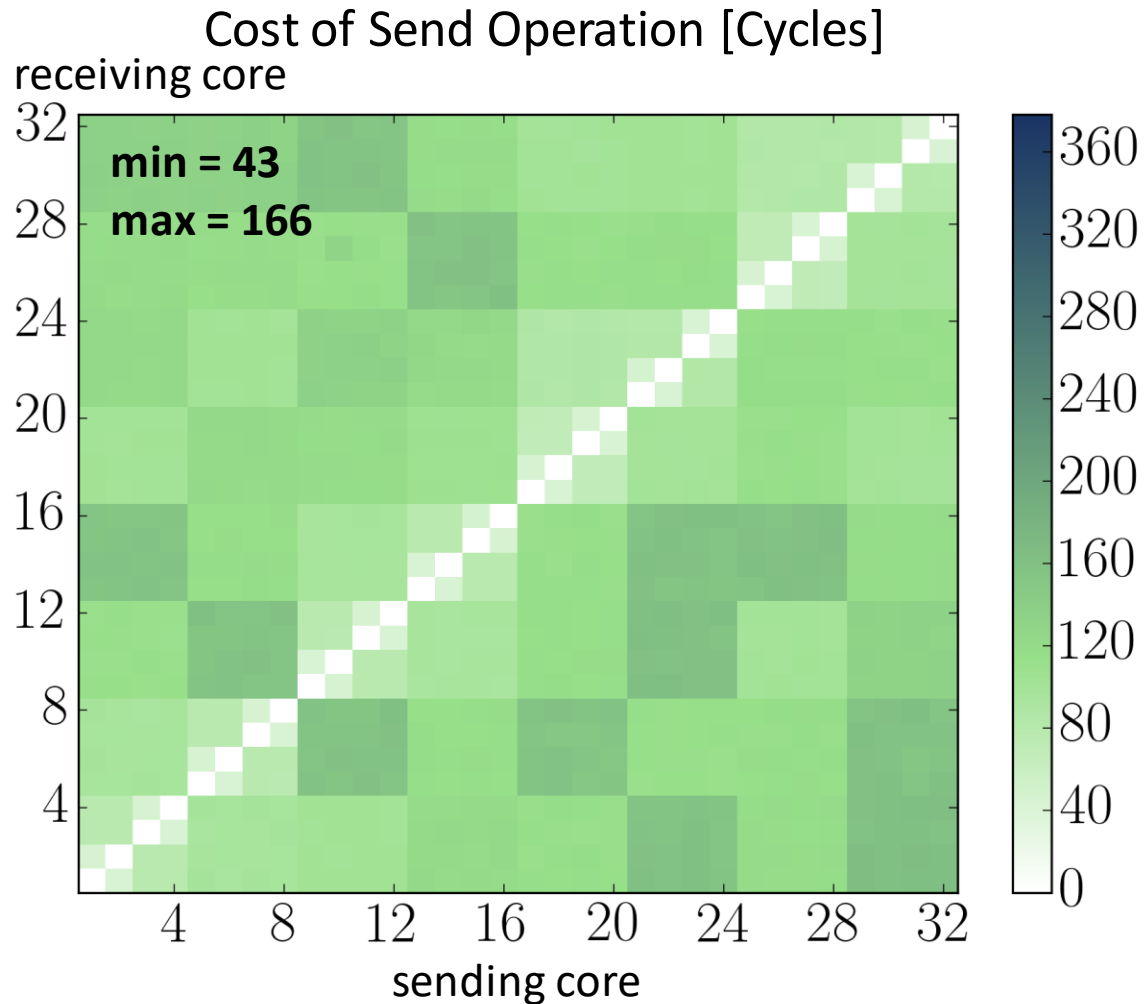
node	0	1	2	3	4	5	6	7
0:	10	16	16	22	16	22	16	22
1:	16	10	22	16	16	22	22	16
2:	16	22	10	16	16	16	16	16
3:	22	16	16	10	16	16	22	22
4:	16	16	16	16	10	16	16	22
5:	22	22	16	16	16	10	22	16
6:	16	22	16	22	16	22	10	16
7:	22	16	16	22	22	16	16	10

2

4

Complement with microbenchmarks: pairwise send and receive

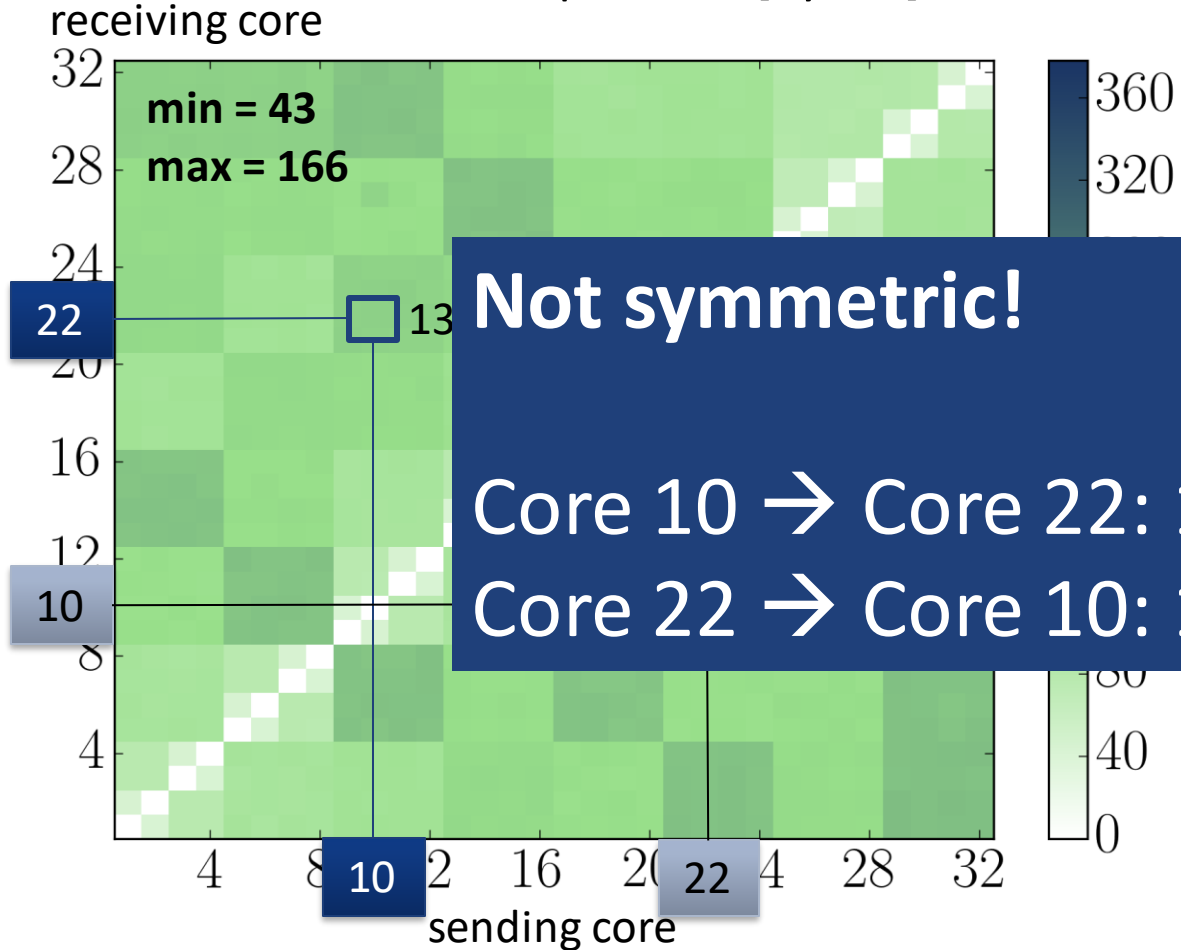
AMD Interlagos 4x4x2



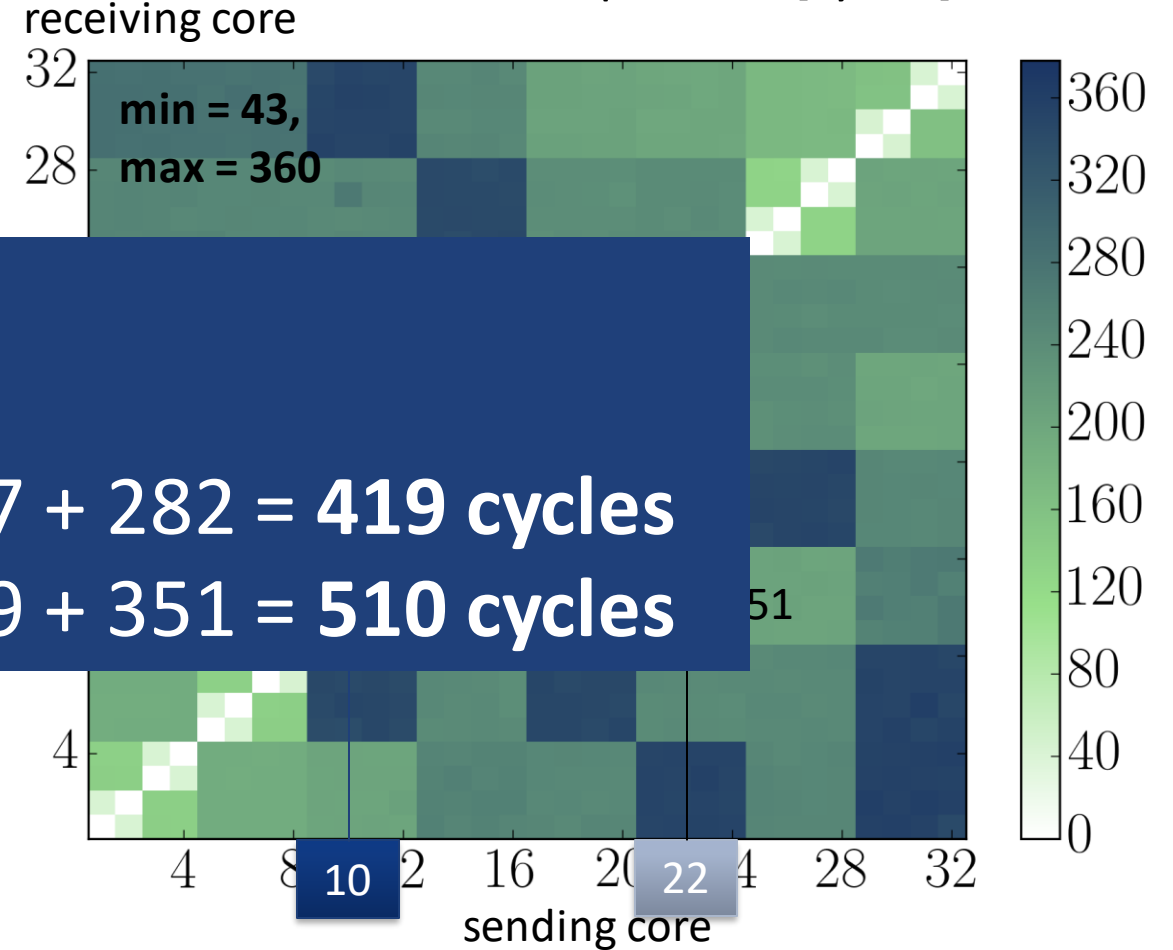
Complement with microbenchmarks: pairwise send and receive

AMD Interlagos 4x4x2

Cost of Send Operation [Cycles]



Cost of Receive Operation [cycles]



Smelt

Using Smelt for group communication

```
#include <smelt/smelt.h>

void main() {

    smelt_init();

    smelt_topology_create();

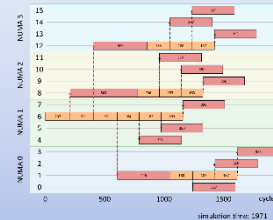
    smelt_broadcast(msg);

}
```

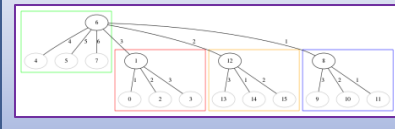
Program

Smelt runtime
smelt_topology_create() {

Smelt Algorithm



Tree topology



}

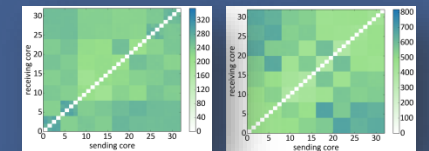
Hardware Discovery

Cores
NUMA nodes

lstopo; /proc/cpu

Measurements

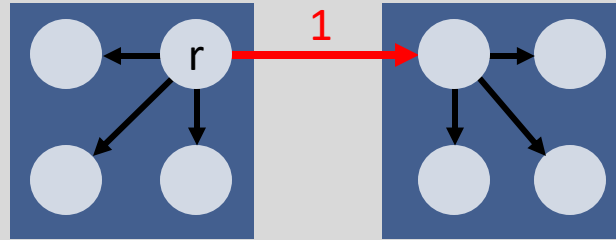
Micro-benchmarks



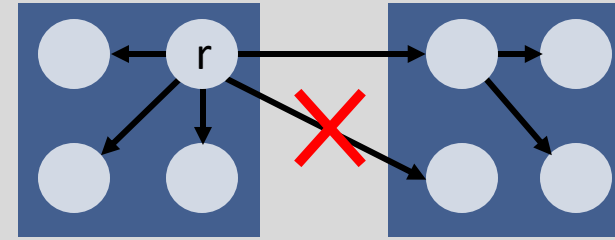
Multicore Model

Smelt's tree generator heuristics

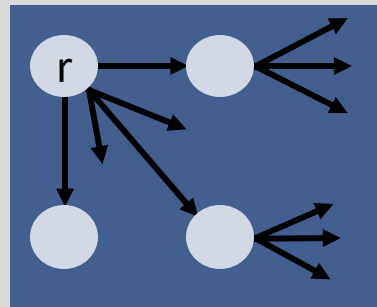
Remote Cores First



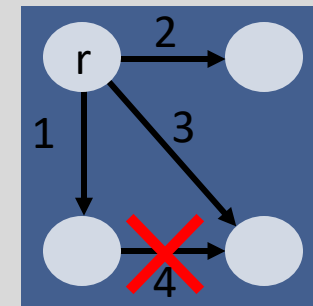
Avoid Expensive Communication



Maximize Parallelism



No redundancy



NUMA 3

15
14
13
12

NUMA 2

11
10
9
8

NUMA 1

7
6
5
4

NUMA 0

3
2
1
0

send

receive

Broadcast on
AMD Shanghai 4x4x1

0

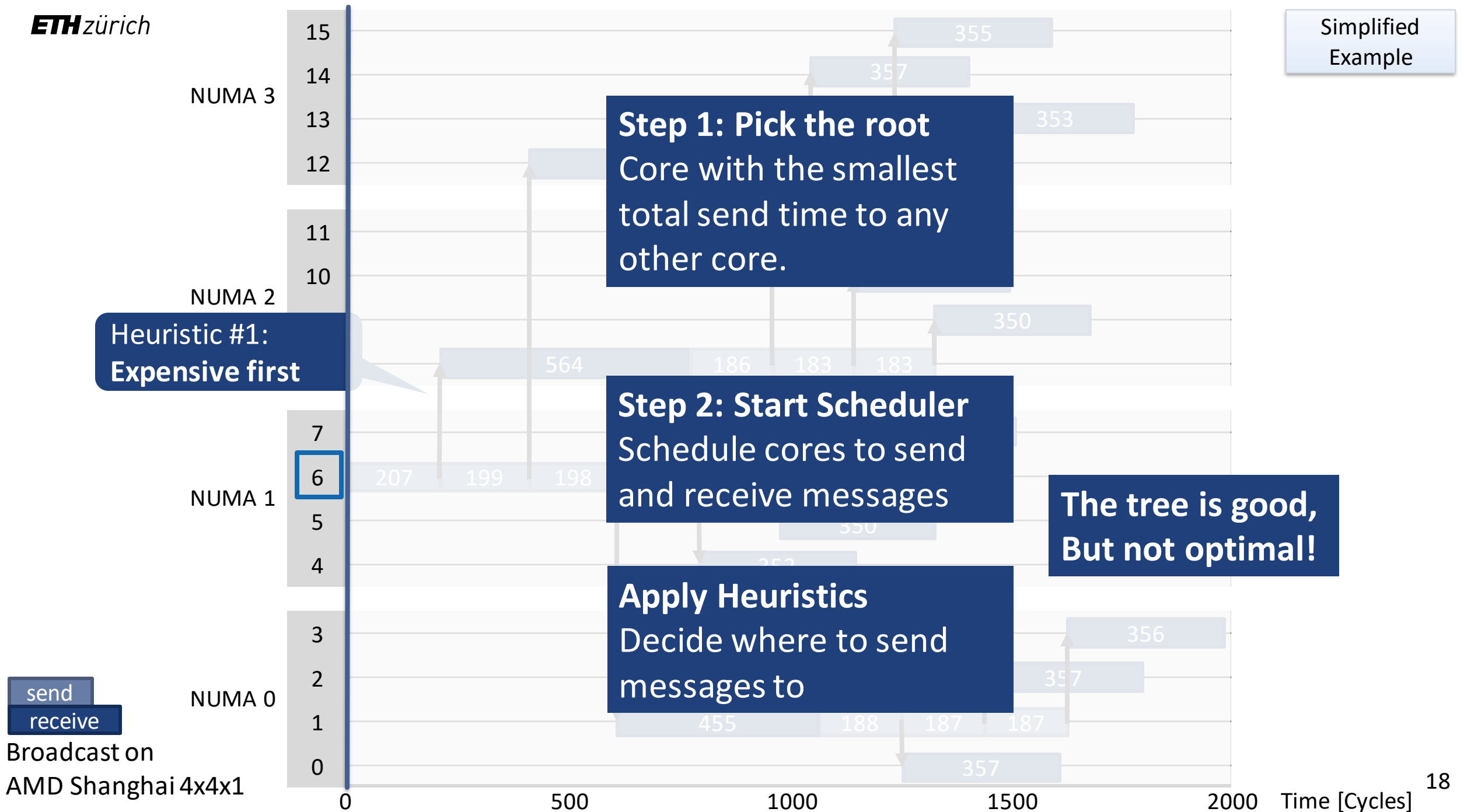
500

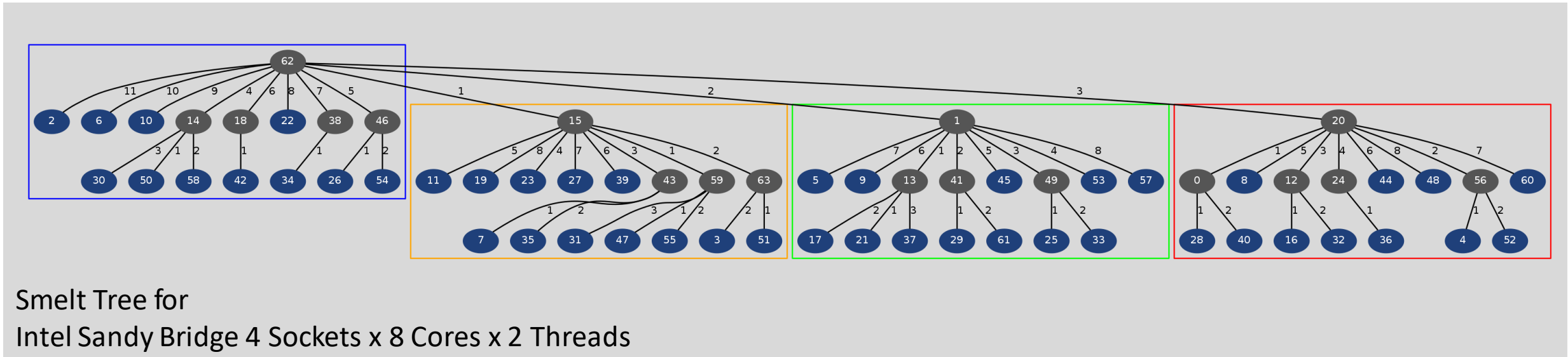
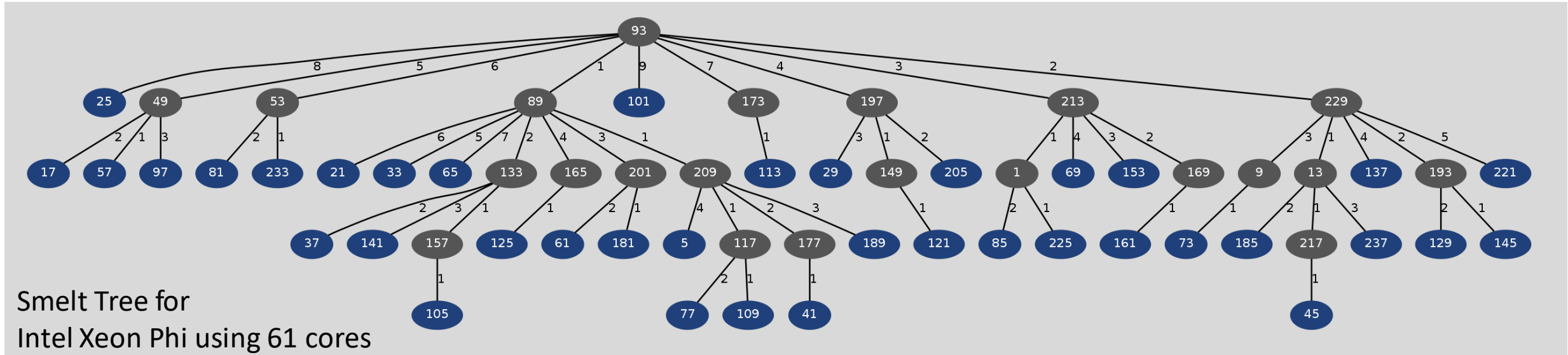
1000

1500

2000

Time [Cycles]





Evaluation Testbed

Intel

Architecture	Sockets	Cores / Socket	Threads / Core
Ivy Bridge	2	10	2
Nehalem	4	8	2
Knights Corner	1	61	4
Sandy Bridge	4	8	2
Sandy Bridge	2	10	2
Bloomfield	2	4	2

AMD

Architecture	Sockets	Cores / Socket	Threads / Core
Magny Cours	4	12	1
Barcelona	8	4	1
Shanghai	4	4	1
Interlagos	4	4	2
Istanbul	4	6	1

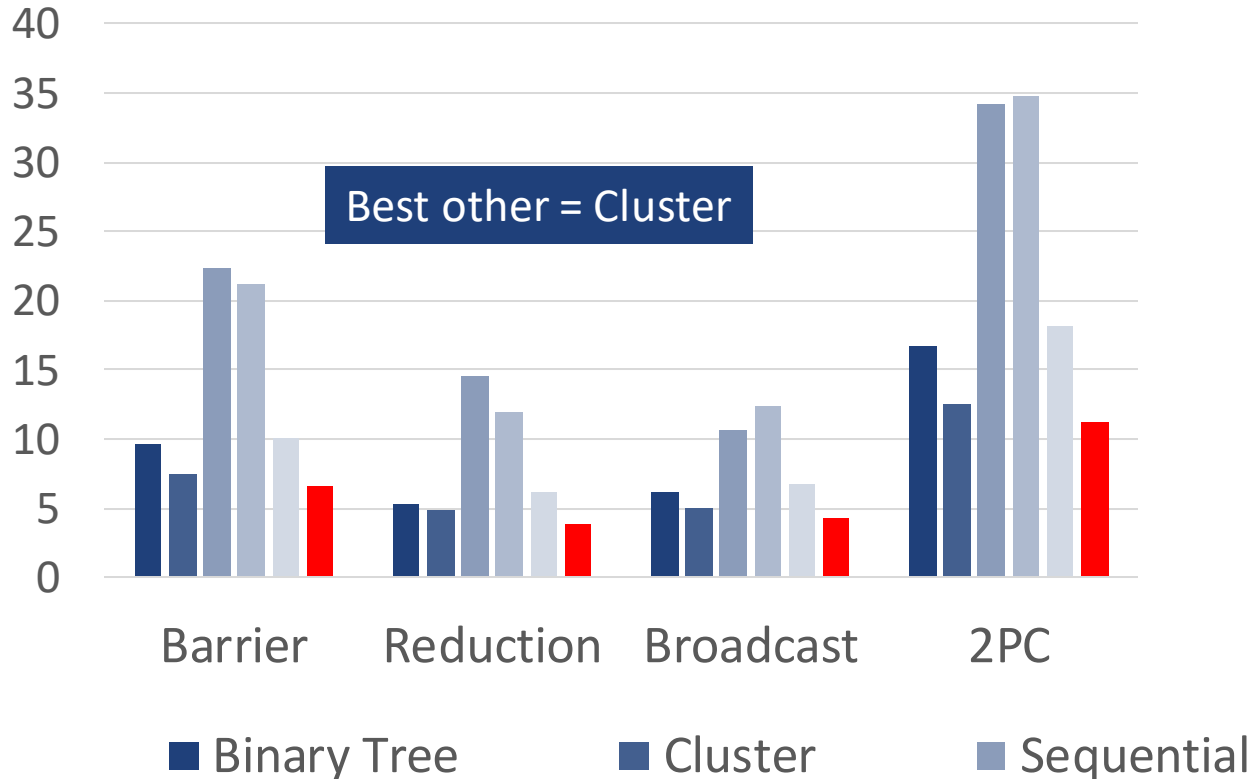
Full set of results online.

<http://machinedb.systems.ethz.ch>

Smelt produces good trees across architectures

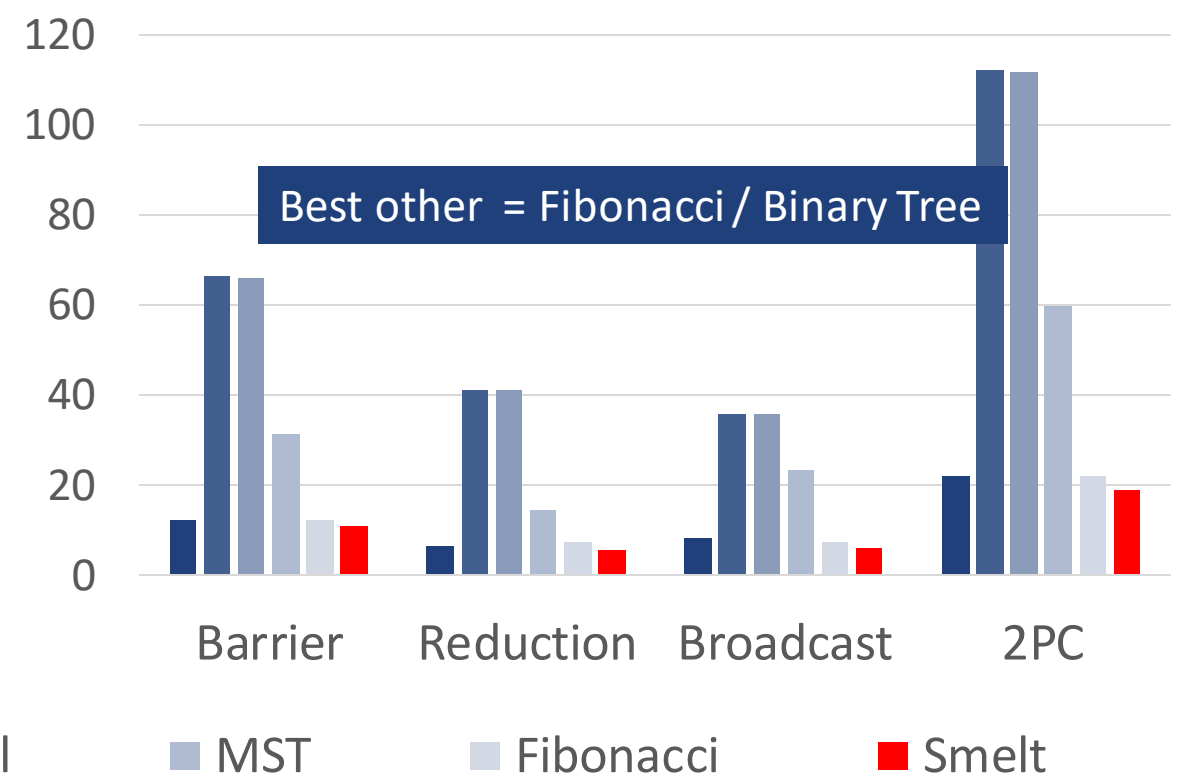
AMD Interlagos (4 Socket x 4 Threads)

Execution Time [kCycles]

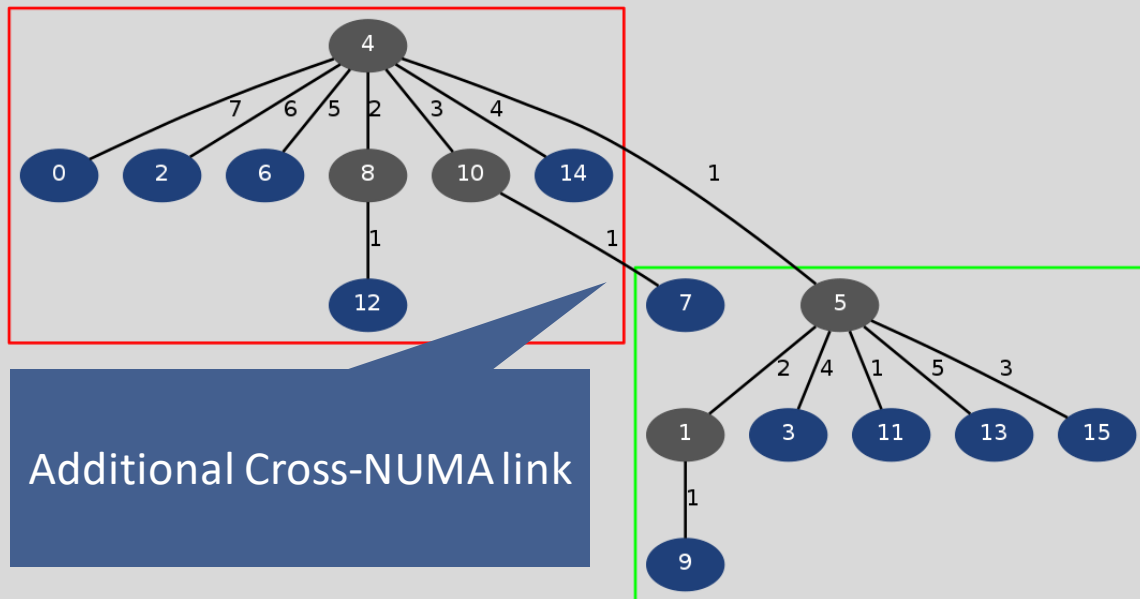


Intel Xeon Phi (61 Threads)

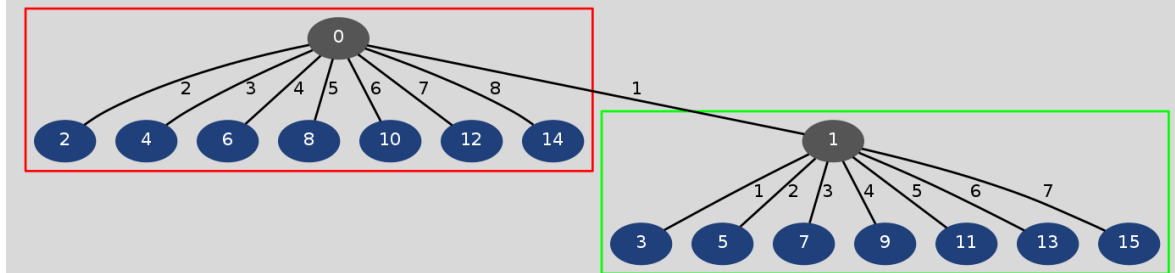
Execution Time [kCycles]



Fast broadcast trees are good for reductions in most cases



Smelt Tree on Intel Bloomfield 2x4x2

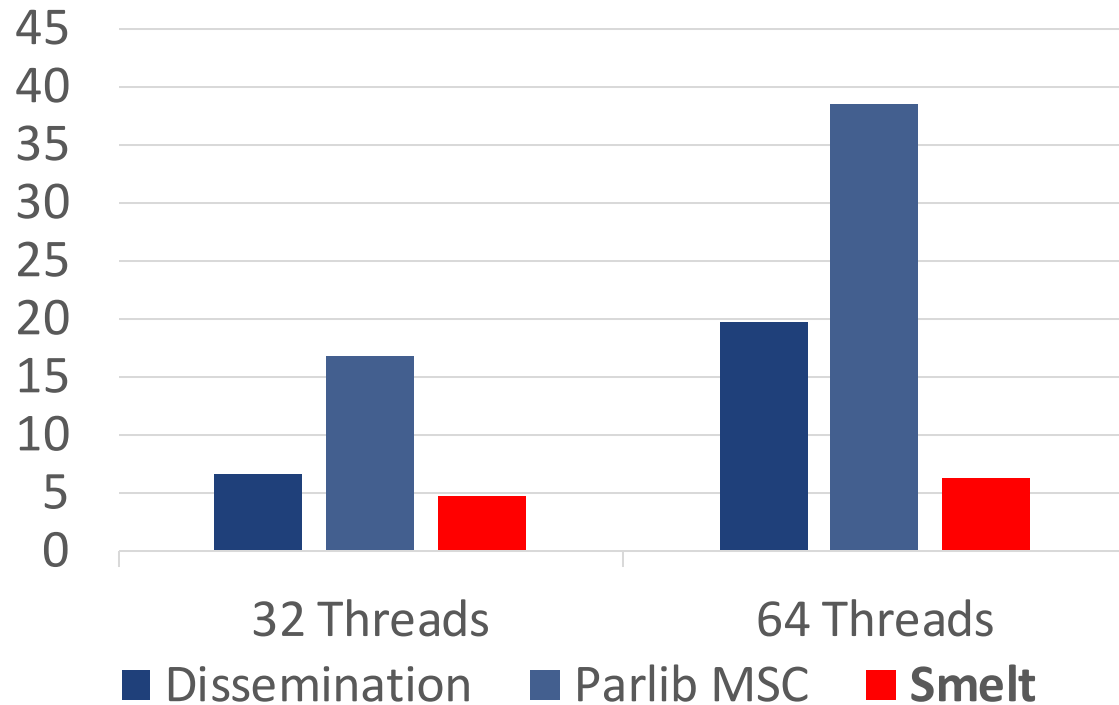


Cluster Topology on Intel Bloomfield 2x4x2

Smelt provides simple and fast barriers

Barrier Benchmark on Intel Sandy Bridge 4x8x2

Execution Time [kCycles]



Barriers based on reduction and broadcast

```
void smelt_barrier(void) {  
    smelt_reduce();  
    smelt_broadcast();  
}
```

Simple barrier implementation

OpenMP: EPCC OpenMP Benchmark Collection

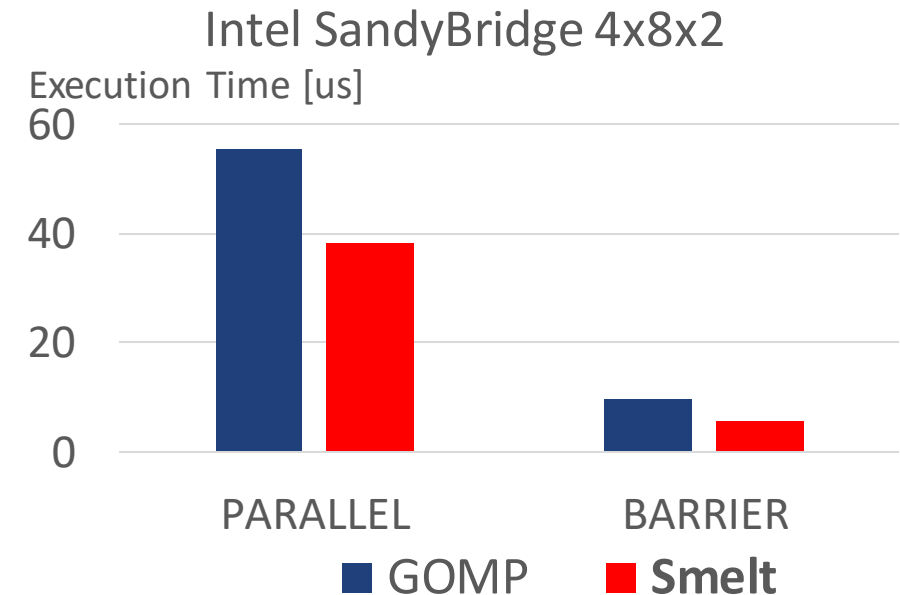
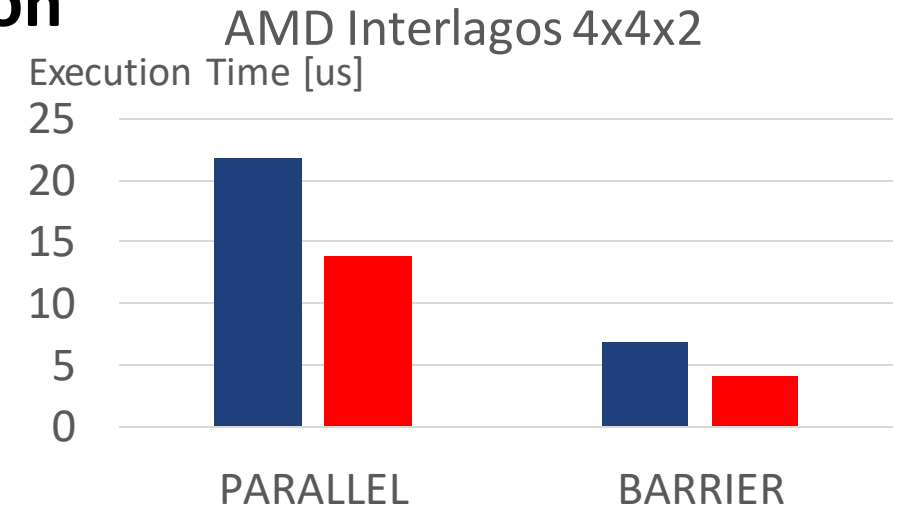
```
/* epcc openmp barrier benchmark */  
void testbar() {  
    int j;  
    #pragma omp parallel private(j)  
    {  
        for (j = 0; j < innerreps; j++) {  
            delay(delaylength);  
            #pragma omp barrier  
        }  
    }  
}
```

Explicit barrier

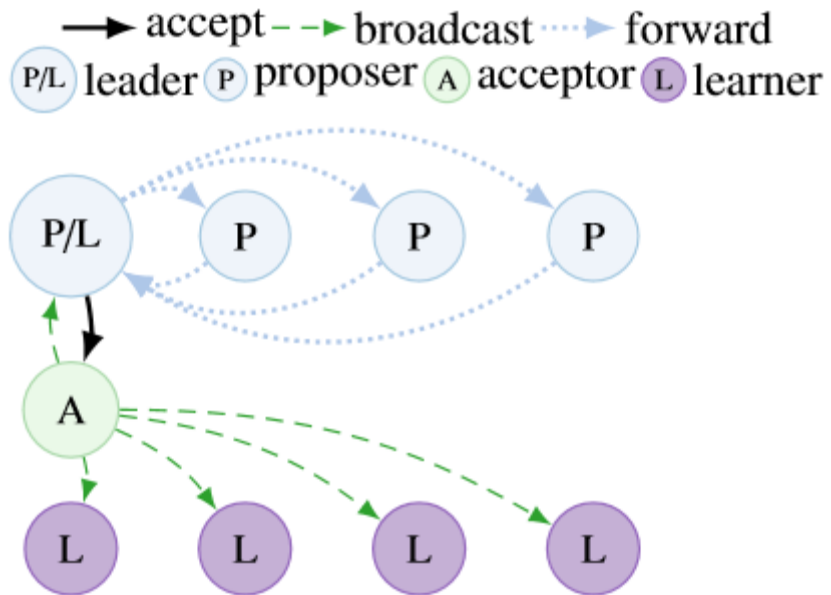
Implicit barrier at the end of parallel block

Replaced GOMP barrier with Smelt

→ Remaining results on the website



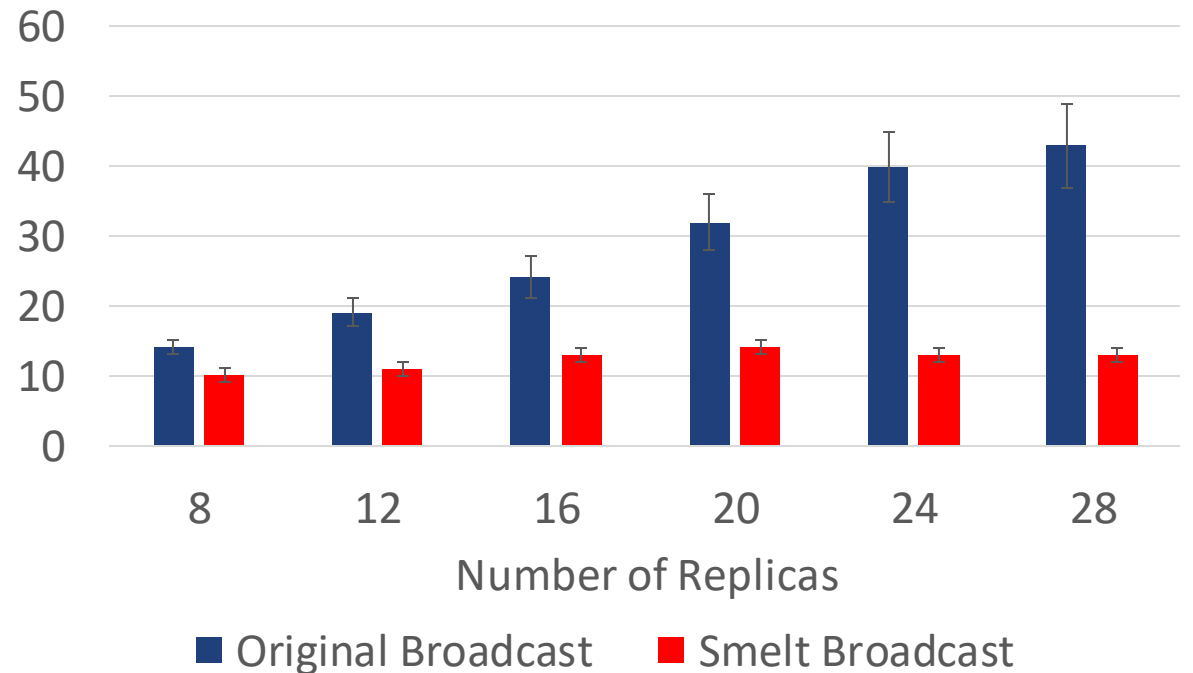
Agreement Protocols: 1Paxos



4 clients to generate load
N replicas executing 1Paxos

1Paxos Benchmark on AMD Interlagos 4x4x2

Execution Time [kCycles]



Summary

- Broadcasts and reductions are **central building blocks**
- **No globally optimal** tree topology
- Information from hardware discovery is **not sufficient**
- Smelt's produces **good trees**

Talk to us at
the first poster
session

